Topology and Geometry of Low-dimensional Manifolds at Nara Women's University October 31, 2018

Outline

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Yamashita's output of Rileyslice

* There is a practical "algorithm" to check if a given two-parabolic group is discrete or not !! [4/18

Demonstration 1



Goal: Replace the "main cusp" with cone singularity.



<u>Rem</u> The double covering space of S³ branched along S(p.g) is the lens space L(p.g).





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Idea of proof — based on Jorgensen's argument A Jorgensen characterized the combinatorial structures of the Ford domains for punctured torus groups.



From one-cone torus to two-bridge cone manifolds

Demonstration 2

Thm [L-S] $QF(M) \cong Teich(\partial M)$



Ford domain for cone hyp str

Suppose
$$\exists CCM$$

s.t. $C \cong glue(I)$

The Ford domain w.r.t. C is





Partial results

Thm [[A, 2015] The Ford domain for any "fuchsian" structure on T20 XR has "good" comb. str. Thm 2 [A, 2018] · The Ford domain for any "thin" structure on TooxIR has "good" comb. str. . The space of thin str is parametrized by the hyp str of 2 (compact convex core). Explosion of isometric hemispheres

In Demo 2,
$$r(ison hemi) \longrightarrow \infty$$
 as $d \longrightarrow 2\pi - 0$.

- 0=0 (cusp case)
 Explosion occurs at d=2TL, where
 we obtain the complete hyp str of 2-br knot.
- 0 > 0 (our new case)
 Explosion occurs before α = 2π
 ⇒ We cannot reach the desired strongly via Ford domains

Employing Dirichlet domain Series-Tan-Yamashita studied the extended Riley slices, where Dirichlet domain is employed.



Demonstration 3

Dirichlet domain

The Dirichlet domain w.r.t. P is

$$M - \left\{ x \mid \overset{\exists 2}{\Rightarrow} \text{ shortest paths} \right\}$$

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Key Lemma for Prop





Choice of base points
Schematic sectional views of Ford/Dirichlet domains:
$\Sigma_0 = \Sigma$ $\longrightarrow \Sigma_0 = \Sigma$ \longrightarrow Σ_0
$d < 2\pi - 0$ $d = 2\pi - 0$ $2\pi - 0 < d \le 2\pi$
Symbolize the axes of "elliptic generators"
of slope 1, 8 for K(p.8).
Conj VpEZo, the Dirichlet domain w.r.t.p
has good comb str.

Demonstration 4

Some comments

- The cone hyp str for the figure 8 knot in Demo 4 are studied by Hilden-Lozano-Montesinos in detail by using Dirichlet domains w.r.t. other base points.
- Another family for bundle over S' is studied by Jorgensen, and Heusener – - Porti - Suárez.

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 The families for figure 8 knot are also studied by Thurston.